Your Censor is My Censor: Weaponizing Censorship Infrastructure for Availability Attacks

Kevin Bock Pranav Bharadwaj



Idwaj Jasraj Singh Dave Levin

UNIVERSITY OF MARYLAND









Web browser







Web browser







Web browser



























kittens













































































Types of residual censorship Categorized by what information censor remembers





Types of residual censorship Categorized by what information censor remembers



4-tuple

3-tuple

2-tuple



Source Destination (IP, port, IP, port) (IP, IP, port) (IP, IP)

More aggressive



Censorship infrastructure



Censorship infrastructure





Censorship infrastructure can be weaponized



















































STOP

















STOP






















puppies











puppies













Attackers can restrict benign communication from crossing the censors' borders





36



How can we evaluate ethically?

Limitations of residual censorship?



How can we

Experiments in Iran, China, and Kazahkstan

Find differences in implementation and duration

Ethical experiments with SP³

Attacked ourselves from a dozen vantage points

Current state of residual censorship?

How can it be weaponized?

Experiments in Iran, China, and Kazahkstan

Find differences in implementation and duration

Current state of residual censorship?

How can it be weaponized?

State of residual censorship

Diversity of censors



Diversity of protocols

SNI ESNI DNS SMTP Other















Types of residual censorship 3-tuple or 4-tuple









Duration of residual censorship How long does blocking last?







Is residual censorship bidirectional? Does it affect traffic entering the country?



ESNI SMTP DNS Other





In all cases, censor tracks traffic direction

State of residual censorship



Residual censorship is implemented differently around the world

Bi-directional, but direction matters

- Different censorship mechanisms (RSTs vs Null Routing)
- Different types of censorship, even within countries

Experiments in Iran, India, China, and Kazahkstan

Find differences in implementation and duration

Ethical experiments with SP³

Attacked ourselves from a dozen vantage points

Current state of residual censorship?

How can it be weaponized?

Attacked ourselves from a dozen vantage points

Current state of residual censorship?

How can it be weaponized?

Ethical experiments with SP³





Weaponizing Middleboxes for **TCP Reflected Amplification** to appear in USENIX Security later this summer



Weaponizing Middleboxes for TCP Reflected Amplification to appear in USENIX Security later this summer

Packet sequences

SYN with Request

PSH

PSH+ACK

SYN; PSH

SYN; PSH+ACK



Weaponizing Middleboxes for TCP Reflected Amplification to appear in USENIX Security later this summer

Censorship can be triggered without a proper 3-way handshake

Packet sequences

SYN with Request

PSH

PSH+ACK

SYN; PSH

SYN; PSH+ACK



Weaponizing Middleboxes for TCP Reflected Amplification to appear in USENIX Security later this summer

Censorship can be triggered without a proper 3-way handshake

Packet sequences

SYN with Request

PSH

PSH+ACK

SYN; PSH

SYN; PSH+ACK





Only between hosts we control

Full control over packets we send

Ethical evaluation

- Attack ourselves ethically without affecting other hosts Need to spoof traffic to or from a censored regime

 - Solution: SP³

















Consent to receive source-spoofed packets



















Forbidden sequence

















Innocuous request

Around the world







Innocuous request

Censored Regimes







"Attacker" (SP³)



"Victims" (Our clients)

Tested from 16 external vantage points









Tested from 16 external vantage points





		Destination Location							
Victim Location				Beijing 1	Beijing 2				
Australia	Sydney	HIIP HIIPS	HIIP HIIPS	HIIP ESNI	HIIP ESINI				
China	Beijing 1 Beijing 2								
India	Mumbai								
	Bangalore 1								
	Bangalore 2								
Iran	Tehran								
Ireland	Dublin 1								
	Dublin 2								
Japan	Tokyo								
Kazakhstan	Qaraghandy								
Russia	Khabarovsk								
UAE	Dubai 1								
	Dubai 2								
USA	Colorado								
	owa								
	Virginia								





Victim Location		Kazakhstan		Iran	Beijing 1		Beijing 2		
		HTTP	HTTPS	HTTP	HTTPS	HTTP	ESNI	HTTP	ESNI
Australia	Sydney				\checkmark	50%	10%	55%	\checkmark
China	Beijing 1	×	\checkmark		\checkmark	N/A	N/A	N/A	N/A
	Beijing 2	×			\checkmark	N/A	N/A	N/A	N/A
India	Mumbai	×	\checkmark		\checkmark	×	×		30%
	Bangalore 1				\checkmark	50%	10%		
	Bangalore 2		\checkmark		\checkmark	25%	10%		\checkmark
Iran	Tehran		\checkmark	N/A	N/A		50%	75%	\checkmark
Ireland	Dublin 1	×	\checkmark		\checkmark	×	×	×	5%
	Dublin 2	×	\checkmark			50%	\mathbf{X}	×	\mathbf{X}
Japan	Tokyo		\checkmark		\checkmark	25%	×	×	\checkmark
Kazakhstan	Qaraghandy	N/A	N/A		\checkmark	50%	×	20%	×
Russia	Khabarovsk		\checkmark		\checkmark		×		×
UAE	Dubai 1	×	\checkmark		\checkmark	85%	\mathbf{X}	95%	×
	Dubai 2	×			\checkmark	×	10%	×	50%
USA	Colorado		\checkmark		\checkmark	×	\mathbf{X}		×
	lowa	×				×		×	60%
	Virginia					85%		55%	$\mathbf{\mathbf{X}}$

Results

Destination Location







Source-spoofed traceroute from both to compare network path

Why does it fail?











Source-spoofed traceroute from both to compare network path

Why does it fail?









Why does it fail?

Source-spoofed traceroute from both to compare network path















Why does it fail?







......












Why does it fail?







Why does it fail?









Why does it fail?









Depends on type of residual censorship

Sustaining the attack Goal: block client IP to server IP:port









Attacker can't guess source port

4-tuple (IP, port, IP, port) 3-tuple (IP, IP, port)





Sustaining the attack Goal: block client IP to server IP:port





Attacker can re-trigger from all 65,535 src ports









Attacker can re-trigger from all 65,535 src ports

















Trigger packets × 65,535

Duration







Speed
requiredTrigger packets× 65,535
× 65,535I45 bytes × 65,535
= 120 seconds= 634 kbps















Speed = Trigger packets required Duration

Weak attacker can launch this attack effectively





China



Sustaining the Attack Victim helps sustain the attack

 $(\bullet \bullet)$









 $(\bullet \bullet)$









Innocuous request





Residual timer resets if the victim sends data

Victim retransmissions unknowingly sustain the attack on themselves







Innocuous request

Can the server detect this?







Attacker can limit TTL of packets to reach censor, but not server

Can the server detect this?



Attacker can limit TTL of packets to reach censor, but not server



Attacker must have a vantage point:

(1) Without egress filtering

3

4) Censor can be triggered statelessly

Attack Limitations

- 2) Shares a similar enough path with their victim
 - Traffic crosses a censor (with residual censorship)

Surprisingly high number of shared network paths

What can be done?

Abolish 3-tuple residual censorship

Properly track 3-way handshake

- Some mitigations available to censorship infrastructure:

 - Null routing should track sequence numbers

Unfortunately, no good countermeasures available to victims

Other details in the paper





Reliability Experiments

Studied the reliability of residual censorship

ESNI Weaponization

Attack

Breadth



Examine which ports are affected

Details on how to weaponize ESNI in China

Analysis of other countries that might be affected







censorship.ai Code and website

Weaponizing Middleboxes

Censors can be weaponized to launch availability attacks

Can be done from a weak attacker

Censors pose a threat to the entire Internet